#### Selection from MIT:Computer Architecture and Organization FEL:Architektury počítačů

Version: 1.0

#### **Practical exercise 7**

**ČVUT-FEL in Prague** 

# **Branch Not Taken**

		Branch to Z A B C D Z		
cycle b	cycle b+1	cycle b+2	cycle b+3	cycle b+4
Branch fetched	Branch decoded	Branch decision	PC keeps D (br. not taken)	
	A fetched	A decoded	Aexecuted	A continues
		B fetched	B decoded	B executed
			C fetched	C decoded
				D fetched



flushed if branch is taken

Z fetched

### **Branch Prediction**

Useful for program loops.
A one-bit prediction scheme: a one-bit buffer carries a "history bit" that tells what happened on the last branch instruction

History bit = 1, branch was taken

History bit = 0, branch was not taken



## **Branch Prediction**



### **Branch Prediction for a Loop**



**Execution of Instruction 4** 

h.bit = 0 branch not taken, h.bit = 1 branch taken.

Two-Bit Prediction Buffer
Can improve correct prediction statistics.



#### **Branch Prediction for a Loop**



**Execution of Instruction 4** 

Execu -tion seq.	Old Pred. Buf	Next instr.			New	Predi
		Pred.	I	Act.	Buf	ction
1	10	2	1	2	—11	Good
2	11 🕇	2	2	2	11	Good
3	_ 11 ←	2	3	2	—11	Good
4	11	2	4	2	—11	Good
5	11	2	5	2	—11	Good
6	11 🕇	2	6	2	_11	Good
7	11 🗕	2	7	2	11	Good
8	11 🔶	2	8	2	-11	Good
9	11 🕇	2	9	2	—11	Good
10	11 🕇	2	10	5	10	Bad